

# Conservative Extensions of Structures

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# Enumerations of Structures

In this talk we shall study a relation between structures  $\mathfrak{A}$  and  $\mathfrak{B}$ , possibly with different signatures and  $|\mathfrak{A}| \subseteq |\mathfrak{B}|$ . The structures will have the form  $\mathfrak{A} = (A; P_1, \dots, P_k, =)$ , where  $A$  is an arbitrary countable set.

- ▶ Let  $f$  be an injective mapping onto  $A$  and  $\text{Dom}(f) \subseteq \mathbb{N}$ .
- ▶ For  $X \subseteq A^m$ , we denote the pullback of  $X$  under  $f$  as  $f^{-1}(X) = \{\langle x_1 \dots x_m \rangle : (f(x_1), \dots, f(x_m)) \in X\}$ .
- ▶ The pullback of  $\mathfrak{A}$  under  $f$  is the set  $f^{-1}(\mathfrak{A}) = f^{-1}(P_1) \oplus \dots \oplus f^{-1}(P_k) \oplus f^{-1}(=)$ .

Informally speaking, you can think of  $f^{-1}(\mathfrak{A})$  as the atomic diagram of the copy of  $\mathfrak{A}$  under  $f$ .

# Enumerations of Structures

## Definition

We call an enumeration of  $\mathfrak{A}$  the pair  $\alpha = (f_\alpha, R_\alpha)$ , where

- (i)  $f_\alpha$  is an injective mapping onto  $A$  and  $\text{Dom}(f_\alpha) \subseteq \mathbb{N}$ ,
- (ii)  $R_\alpha \subseteq \mathbb{N}$  and  $f_\alpha^{-1}(\mathfrak{A}) \leq_T R_\alpha$ .

For an enumeration  $\alpha$ , we denote  $\alpha^{(n)} = (f_\alpha, R_\alpha^{(n)})$ .

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## Definition (Comparing two enumerations)

Let  $\mathfrak{A}$  and  $\mathfrak{B}$  be countable structures with  $|\mathfrak{A}| \subseteq |\mathfrak{B}|$ ,  $\alpha = (f_\alpha, R_\alpha)$  and  $\beta = (f_\beta, R_\beta)$  are enumerations of  $\mathfrak{A}$  and  $\mathfrak{B}$  respectively. We denote  $\alpha \leq \beta$  if:

- (i)  $R_\alpha \leq_T R_\beta$ , and
- (ii)  $E(f_\alpha, f_\beta)$  is c.e. in  $R_\beta$ , where

$$E(f_\alpha, f_\beta) = \{\langle x, y \rangle \mid x \in \text{Dom}(f_\alpha) \ \& \ y \in \text{Dom}(f_\beta) \ \& \ f_\alpha(x) = f_\beta(y)\}.$$

# Comparing Structures

## Definition

Let  $\mathfrak{A}$  and  $\mathfrak{B}$  be countable structures, possibly with different signatures and  $|\mathfrak{A}| \subseteq |\mathfrak{B}|$ .

- (i)  $\mathfrak{A} \Rightarrow \mathfrak{B}$ , if for every enumeration  $\beta$  of  $\mathfrak{B}$ , there exists an enumeration  $\alpha$  of  $\mathfrak{A}$ , such that  $\alpha \leq \beta$ .
- (ii)  $\mathfrak{A} \Leftarrow \mathfrak{B}$ , if for every enumeration  $\alpha$  of  $\mathfrak{A}$ , there exists an enumeration  $\beta$  of  $\mathfrak{B}$ , such that  $\beta \leq \alpha$ .
- (iii)  $\mathfrak{A} \Leftrightarrow \mathfrak{B}$ , if  $\mathfrak{A} \Rightarrow \mathfrak{B}$  and  $\mathfrak{A} \Leftarrow \mathfrak{B}$ .

# Comparing Structures

We generalize the previous definition for arbitrary  $k, n \in \mathbb{N}$ .

## Definition

Let  $\mathfrak{A}$  and  $\mathfrak{B}$  be countable structures, possibly with different signatures and  $|\mathfrak{A}| \subseteq |\mathfrak{B}|$ .

- (i)  $\mathfrak{A} \Rightarrow_n^k \mathfrak{B}$ , if for every enumeration  $\beta$  of  $\mathfrak{B}$ , there exists an enumeration  $\alpha$  of  $\mathfrak{A}$ , such that  $\alpha^{(k)} \leq \beta^{(n)}$ .
- (ii)  $\mathfrak{A} \Leftarrow_n^k \mathfrak{B}$ , if for every enumeration  $\alpha$  of  $\mathfrak{A}$ , there exists an enumeration  $\beta$  of  $\mathfrak{B}$ , such that  $\beta^{(k)} \leq \alpha^{(n)}$ .
- (iii)  $\mathfrak{A} \Leftrightarrow_n^k \mathfrak{B}$ , if  $\mathfrak{A} \Rightarrow_n^k \mathfrak{B}$  and  $\mathfrak{A} \Leftarrow_n^k \mathfrak{B}$ .

In this case we say that  $\mathfrak{B}$  is  $(k, n)$ -conservative extension of  $\mathfrak{A}$ .

# Characterization

## Definition (Computationally infinitary formulas)

- ▶  $\Sigma_0^c$  and  $\Pi_0^c$  formulas are quantifier-free formulas with a finite number of conjunctions and disjunctions.
- ▶  $\Sigma_{n+1}^c$  formulas have the form  $\varphi(\bar{x}) \equiv \bigvee_{i \in W_e} \exists \bar{y} \psi_i(\bar{x}, \bar{y})$ , where  $W_e$  is a c.e. set with index  $e$  and  $\psi_i$  are  $\Pi_n^c$  formulas.
- ▶  $\Pi_{n+1}^c$  formulas are negations of  $\Sigma_{n+1}^c$  formulas.

## Definition

We say that a set  $X \subseteq A$  is definable in  $\mathfrak{A}$  with a  $\Sigma_n^c$  formula, denoted with  $X \in \Sigma_n^c(\mathfrak{A})$ , if there exists a tuple of constants  $\bar{c}$  in  $|\mathfrak{A}|$  and a  $\Sigma_n^c$  formula  $\psi(x, \bar{y})$  such that

$$x \in X \iff \mathfrak{A} \models \psi(x, \bar{c}).$$

# Characterization

Theorem (Ash-Knight-Manasse-Slaman, Chisholm)

Let  $\mathfrak{A}$  be a countable structure. For every set  $X \subseteq A$ ,

$X \in \Sigma_{n+1}^c(\mathfrak{A}) \iff (\forall \alpha \text{ enumeration of } \mathfrak{A}) [f^{-1}(X) \text{ is c.e. in } R_\alpha^{(n)}].$



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The following theorem gives the reason for our choice of notation.

**Theorem**

Let  $\mathfrak{A}$  and  $\mathfrak{B}$  be countable structures and  $|\mathfrak{A}| \subseteq |\mathfrak{B}|$ . Then for every  $k, n \in \mathbb{N}$ ,

- (i) if  $\mathfrak{A} \Rightarrow_n^k \mathfrak{B}$ , then  $(\forall X \subseteq A)[X \in \Sigma_{k+1}^c(\mathfrak{A}) \rightarrow X \in \Sigma_{n+1}^c(\mathfrak{B})]$ ;
- (ii) if  $\mathfrak{A} \Leftarrow_n^k \mathfrak{B}$ , then  $(\forall X \subseteq A)[X \in \Sigma_{n+1}^c(\mathfrak{B}) \rightarrow X \in \Sigma_{k+1}^c(\mathfrak{A})]$ ;
- (iii) if  $\mathfrak{A} \Leftrightarrow_n^k \mathfrak{B}$ , then  $(\forall X \subseteq A)[X \in \Sigma_{k+1}^c(\mathfrak{A}) \leftrightarrow X \in \Sigma_{n+1}^c(\mathfrak{B})]$ .

In the general case, we do not have the reverse implications.

## Moschovakis' Extension

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- ▶ We can encode the natural numbers in  $A^*$ . Let  $0^* = \bar{0}$  and  $(n + 1)^* = \Pi(\bar{0}, n^*)$ . We denote them by  $\mathbb{N}^*$ .

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- ▶ Moreover, we can encode finite sequences of elements of  $A^*$ . For every  $t_1, \dots, t_{n+1} \in A^*$ , let  $\Pi_1(t_1) = t_1$  and  $\Pi_{n+1}(t_1, \dots, t_{n+1}) = \Pi(t_1, \Pi_n(t_2, \dots, t_{n+1}))$

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- ▶  $A^*$  resembles  $\mathbb{HIF}(A)$ .

## Moschovakis' Extension

The next step is to define functions for decoding. Let  $L$  and  $R$  be function onto  $A^*$  such that:

$$L(\bar{0}) = R(\bar{0}) = 0^*;$$

$$(\forall t \in A)[L(t) = R(t) = 1^*];$$

$$(\forall s, t \in A^*)[L(\Pi(s, t)) = s \ \& \ R(\Pi(s, t)) = t].$$

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### Definition (Moschovakis' Extension)

Let a structure  $\mathfrak{A}$  be given. For every predicate  $P_i$  of  $\mathfrak{A}$ , we define the unary predicate  $P_i^*$  on  $A^*$  such that:

$$P_i^*(t) \iff (\exists a_1, \dots, a_{n_i} \in A)[t = \Pi_{n_i}(a_1, \dots, a_{n_i}) \ \& \ P_i(a_1, \dots, a_{n_i})].$$

Moschovakis' extension of  $\mathfrak{A}$  is the structures

$$\mathfrak{A}^* = (A^*; A_0, P_1^*, \dots, P_k^*, G_\Pi, G_L, G_R, =),$$

where  $G_\Pi$ ,  $G_L$  and  $G_R$  are the graphs of  $\Pi$ ,  $L$  и  $R$ .



# Characterization

Let  $\mathfrak{A}$  and  $\mathfrak{B}$  be countable structures with  $|\mathfrak{A}| \subseteq |\mathfrak{B}|$  and suppose we use the same coding for all structures.

## Proposition

$\mathfrak{A} \Leftrightarrow_n^k \mathfrak{B}$  if and only if  $\mathfrak{A}^* \Leftrightarrow_n^k \mathfrak{B}^*$ . In particular,  $\mathfrak{A} \Leftrightarrow_n^n \mathfrak{A}^*$ .

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## Theorem

Suppose that  $|\mathfrak{A}^*| \subseteq |\mathfrak{B}|$  and

$$(\forall X \subseteq A^*) [X \in \Sigma_{k+1}^c(\mathfrak{A}^*) \rightarrow X \in \Sigma_{n+1}^c(\mathfrak{B})].$$

Then  $\mathfrak{A} \Rightarrow_n^k \mathfrak{B}$ .

## Corollary

$\mathfrak{A} \Rightarrow_n^k \mathfrak{B}$  iff  $(\forall X \subseteq A^*) [X \in \Sigma_{k+1}^c(\mathfrak{A}^*) \rightarrow X \in \Sigma_{n+1}^c(\mathfrak{B}^*)]$ .

# Applications. Degree Spectra

## Definition

- ▶ The Turing degree spectrum of  $\mathfrak{A}$  is the set  $DS(\mathfrak{A}) = \{d_T(R_\alpha) \mid \alpha \text{ is an enumeration of } \mathfrak{A}\}$ .
- ▶ The  $k$ -th Turing jump spectrum of  $\mathfrak{A}$  is the set  $DS_k(\mathfrak{A}) = \{d_T(R_\alpha^{(k)}) \mid \alpha \text{ is an enumeration of } \mathfrak{A}\}$ .

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## Proposition

Let  $\mathfrak{A}$  and  $\mathfrak{B}$  be countable structures with  $|\mathfrak{A}| \subseteq |\mathfrak{B}|$ .

- ▶ If  $\mathfrak{A} \Rightarrow_n^k \mathfrak{B}$ , then  $DS_n(\mathfrak{B}) \subseteq DS_k(\mathfrak{A})$ ;
- ▶ If  $\mathfrak{A} \Leftarrow_k^n \mathfrak{B}$ , then  $DS_k(\mathfrak{A}) \subseteq DS_n(\mathfrak{B})$ ;
- ▶ If  $\mathfrak{A} \Leftrightarrow_n^k \mathfrak{B}$ , then  $DS_k(\mathfrak{A}) = DS_n(\mathfrak{B})$ ;

In the general case, we do not have the reverse implications.

## Applications. The Jump Structure

Let a structure  $\mathfrak{A}$  be given and  $P$  be a relation in  $\mathfrak{A}$ ,  $\tau \in \text{Fin}(\mathfrak{A})$ .

$$\tau^{-1}(P)(u) \downarrow = 1 \leftrightarrow (\exists x_1, \dots, x_n \in \text{Dom}(\tau))[u = \langle x_1, \dots, x_n \rangle \ \& \ (\tau(x_1), \dots, \tau(x_n)) \in P],$$

$$\tau^{-1}(P)(u) \downarrow = 0 \leftrightarrow (\exists x_1, \dots, x_n \in \text{Dom}(\tau))[u = \langle x_1, \dots, x_n \rangle \ \& \ (\tau(x_1), \dots, \tau(x_n)) \notin P].$$

By  $\tau^{-1}(\mathfrak{A})$  we denote the finite function  $\tau^{-1}(P_1) \oplus \dots \oplus \tau^{-1}(P_k)$ .

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### Definition (Forcing)

$$\blacktriangleright \tau \Vdash_0 F_e(x) \leftrightarrow x \in W_e^{\tau^{-1}(\mathfrak{A})},$$

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- ▶  $\tau \Vdash_0 F_e(x) \leftrightarrow x \in W_e^{\tau^{-1}(\mathfrak{A})}$ ,
- ▶  $\tau \Vdash_{n+1} F_e(x) \leftrightarrow (\exists \delta \in \text{Fin}(\{0, 1\})) [x \in W_e^\delta \ \& \ (\forall z \in \text{Dom}(\delta)) [(\delta(z) = 1 \ \& \ \tau \Vdash_n F_z(z)) \vee (\delta(z) = 0 \ \& \ \tau \Vdash_n \neg F_z(z))]]$ ,
- ▶  $\tau \Vdash_n \neg F_e(x) \leftrightarrow (\forall \rho \in \text{Fin}(A)) [\tau \subseteq \rho \rightarrow \rho \not\Vdash_n F_e(x)]$ .

## Applications. The Jump Structure

We define in some sense an analogue of the Kleene set.

$$K_n^{\mathfrak{A}} = \{\Pi_3(\delta^*, e^*, x^*) \mid (\exists \tau \in \text{Fin}(A))[\delta \subseteq \tau \ \& \ \tau \Vdash_n F_e(x)] \ \& \ e^*, x^* \in \mathbb{N}^*\}$$

### Definition

The  $n$ -th jump structure of  $\mathfrak{A}$  is defined as:

$$\mathfrak{A}^{(0)} = \mathfrak{A} \text{ and } \mathfrak{A}^{(n+1)} = (\mathfrak{A}^*, K_n^{\mathfrak{A}}).$$



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- ▶  $\mathfrak{A} \Leftrightarrow_0^n \mathfrak{A}^{(n)}$ . Thus,  $DS_n(\mathfrak{A}) = DS(\mathfrak{A}^{(n)})$ .
- ▶ Moreover,  $(\forall X \subseteq A)[X \in \Sigma_{n+2}^c(\mathfrak{A}) \leftrightarrow X \in \Sigma_{n+1}^c(\mathfrak{A}')] ;$

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- ▶ Moreover,  $(\forall X \subseteq A)[X \in \Sigma_{n+2}^c(\mathfrak{A}) \leftrightarrow X \in \Sigma_{n+1}^c(\mathfrak{A}')] ;$
- ▶  $K_n^{\mathfrak{A}} \in \Sigma_{n+1}^c(\mathfrak{A}^*)$  and  $A^* \setminus K_n^{\mathfrak{A}} \notin \Sigma_{n+1}^c(\mathfrak{A}^*)$ .
- ▶  $\mathfrak{A}^{(n)} \Rightarrow_0 \mathfrak{A}^{(n+1)}$  and  $\mathfrak{A}^{(n)} \not\Rightarrow_0 \mathfrak{A}^{(n+1)}$ .

Therefore, the sequence  $\{\mathfrak{A}, \mathfrak{A}^{(1)}, \mathfrak{A}^{(2)}, \dots, \mathfrak{A}^{(n)}, \dots\}$  is strictly increasing under the relation  $\Rightarrow_0$ .

## Applications. Marker's extension

For  $\mathfrak{A} = (A; P_1, \dots, P_k, =)$ , the Marker's extension of  $\mathfrak{A}$  is a structure of the form

$$\mathfrak{A}^{\exists\forall} = (A \cup \bigcup_{i=1}^k X_i \cup \bigcup_{i=1}^k Y_i; P_1^{\exists\forall}, \dots, P_k^{\exists\forall}, X_1, \dots, X_k, Y_1, \dots, Y_k, =).$$

Among others, it has the following property:

$$\mathfrak{A} \models P_i(\bar{a}) \text{ if and only if } \mathfrak{A}^{\exists\forall} \models (\exists x \in X_i)(\forall y \in Y_i)[P_i^{\exists\forall}(\bar{a}, x, y)]$$

### Theorem (A. Soskova - I. Soskov)

Let  $\mathfrak{A}$  and  $\mathfrak{B}$  be structures with  $DS(\mathfrak{A}) \subseteq DS_1(\mathfrak{B})$ .

Then there exists a structure  $\mathfrak{C}$  such that

$DS(\mathfrak{C}) \subseteq DS(\mathfrak{B})$  and  $DS_1(\mathfrak{C}) = DS(\mathfrak{A})$ .

## Applications. Marker's Extension

We give an analogue of their result. Let  $\mathcal{O}_A = (A, =)$ .

### Proposition (Jump Inversion)

Let  $\mathfrak{A}$  be a countable structure and  $\mathcal{O}_A \Rightarrow_0^k \mathfrak{A}$  for some  $k \in \mathbb{N}$ . There exists a countable structure  $\mathfrak{B}$  such that  $\mathfrak{A} \Leftrightarrow_0^0 \mathfrak{B}^{(k)}$ . In other words,

$$(\forall X \subseteq A)[X \in \Sigma_1^c(\mathfrak{A}) \leftrightarrow X \in \Sigma_1^c(\mathfrak{B}^{(k)})].$$

It is still open whether:

- ▶ We can relativize it, i.e. replace  $\mathcal{O}_A \Rightarrow_0^k \mathfrak{A}$  with  $\mathfrak{C} \Rightarrow_0^k \mathfrak{A}$  for arbitrary  $\mathfrak{C}$ .
- ▶ Lift the result to infinite ordinals.

БЛАГОДАРЯ!  
(Thank you!)