

Fine Hierarchy of ω -Regular k -Partitions

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In [W79] K. Wagner gave in a sense the finest possible topological classification of regular ω -languages (i.e., of the subsets of X^ω for a finite alphabet X recognized by finite automata) known as the Wagner hierarchy. In particular, he completely described the (quotient structure of the) preorder $(\mathcal{R}; \leq_{CA})$ formed by the class \mathcal{R} of regular subsets of X^ω and the reducibility by functions continuous in the Cantor topology on X^ω (note that in descriptive set theory the CA -reducibility is widely known as the Wadge reducibility).

In [S94, S95, S98] the Wagner hierarchy of regular ω -languages was related to the Wadge hierarchy and to the author's fine hierarchy [S95a]. This provided new proofs of results in [W79] and yielded some new results on the Wagner hierarchy. See also alternative algebraic approaches [CP97, CP99, DR06] and [CD09]. The aim of this paper is to generalize this theory from the case of regular ω -regular languages to the case of regular k -partitions of X^ω , i.e. k -tuples (A_0, \dots, A_{k-1}) of pairwise disjoint regular sets satisfying $A_0 \cup \dots \cup A_{k-1} = X^\omega$. Note that the ω -languages are in a bijective correspondence with 2-partitions of X^ω .

- 1) The structure $(\mathcal{R}; \leq_{CA})$ is almost well-ordered with the order type ω^ω , i.e. there are $A_\alpha \in \mathcal{R}$, $\alpha < \omega^\omega$, such that $A_\alpha <_{CA} A_\alpha \oplus \bar{A}_\alpha <_{CA} A_\beta$ for $\alpha < \beta < \omega^\omega$ and any regular set is CA -equivalent to one of the sets $A_\alpha, \bar{A}_\alpha, A_\alpha \oplus \bar{A}_\alpha$ ($\alpha < \omega^\omega$).
- 2) The CA -reducibility coincides on \mathcal{R} with the DA -reducibility, i.e. the reducibility by functions computed by deterministic asynchronous finite transducers, and \mathcal{R} is closed under the DA -reducibility.
- 3) Any level $\mathcal{R}_\alpha = \{C \mid C \leq_{DA} A_\alpha\}$ of the Wagner hierarchy is decidable.

A Muller k -acceptor is a pair (\mathcal{A}, c) where \mathcal{A} is an automaton and $c : C_{\mathcal{A}} \rightarrow k$ is a k -partition of $C_{\mathcal{A}} = \{f_{\mathcal{A}}(\xi) \mid \xi \in X^{\omega}\}$ where $f_{\mathcal{A}}(\xi)$ is the set of states which occur infinitely often in the sequence $f(i, \xi) \in Q^{\omega}$. Note that in this paper we consider only deterministic finite automata. Such a k -acceptor recognizes the k -partition $L(\mathcal{A}, c) = c \circ f_{\mathcal{A}}$ where $f_{\mathcal{A}} : X^{\omega} \rightarrow C_{\mathcal{A}}$ is the map defined above. We have the following characterization of the ω -regular partitions.

Proposition

A partition $L : X^{\omega} \rightarrow k$ is regular iff it is recognized by a Muller k -acceptor.

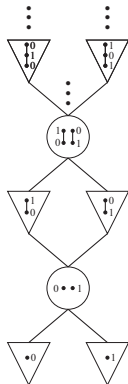
Let $(Q; \leq)$ be a poset. A Q -poset is a triple (P, \leq, c) consisting of a finite nonempty poset $(P; \leq)$, $P \subseteq \omega$, and a labeling $c : P \rightarrow Q$. A *morphism* $f : (P, \leq, c) \rightarrow (P', \leq', c')$ of Q -posets is a monotone function $f : (P; \leq) \rightarrow (P'; \leq')$ satisfying $\forall x \in P (c(x) \leq c'(f(x)))$. Let \mathcal{P}_Q , \mathcal{F}_Q and \mathcal{T}_Q denote the sets of all finite Q -posets, Q -forests and Q -trees, respectively.

The h -preorder \leq_h on \mathcal{P}_Q is defined as follows: $P \leq_h P'$, if there is a morphism from P to P' . Note that for the particular case $Q = \bar{k}$ of the antichain with k elements we obtain the preorders \mathcal{P}_k , \mathcal{F}_k and \mathcal{T}_k .

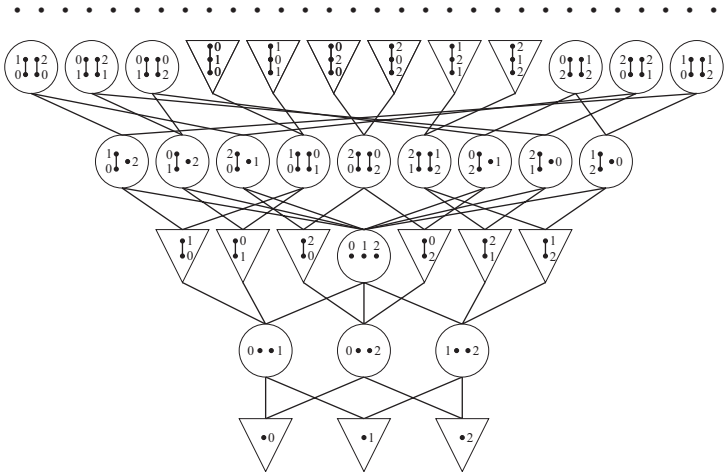
It is well known that if Q is a wqo then $(\mathcal{F}_Q; \leq_h)$ and $(\mathcal{T}_Q; \leq_h)$ are wqo's. Obviously, $P \subseteq Q$ implies $\mathbb{F}_P \subseteq \mathbb{F}_Q$, and $P \sqsubseteq Q$ (i.e., P is an initial segment of Q) implies $\mathbb{F}_P \sqsubseteq \mathbb{F}_Q$.

Define the sequence $\{\mathcal{F}_k(n)\}_{n < \omega}$ of preorders by induction on n as follows: $\mathcal{F}_k(0) = \bar{k}$ and $\mathcal{F}_k(n+1) = \mathcal{F}_{\mathcal{F}_k(n)}$. Identifying the elements $i < k$ of \bar{k} with the corresponding minimal elements $s(i)$ of $\mathcal{F}_k(1)$, we may think that $\mathcal{F}_k(0) \sqsubseteq \mathcal{F}_k(1)$, hence $\mathcal{F}_k(n) \sqsubseteq \mathcal{F}_k(n+1)$ for each $n < \omega$ and $\mathcal{F}_k(\omega) = \bigcup_{n < \omega} \mathcal{F}_k(n)$ is a wqo.

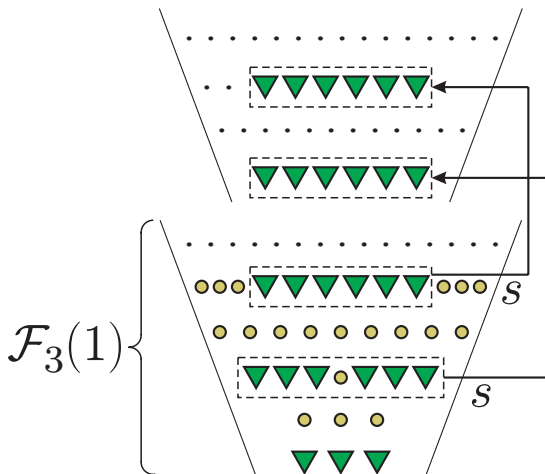
The preorders $\mathcal{F}_k(\omega)$, $\mathcal{T}_k(\omega)$ and the set $\mathcal{T}_k^{\sqcup}(\omega)$ of finite joins of elements in $\mathcal{T}_k(\omega)$, play an important role in the study of the FH of k -partitions because they provide convenient naming systems for the levels of this hierarchy (similar to the previous work where \mathcal{F}_k and \mathcal{T}_k were used to name the levels of the DH of k -partitions). Note that $\mathcal{F}_k(1) = \mathcal{F}_k$ and $\mathcal{T}_k(1) = \mathcal{T}_k$. For the FH of ω -regular k -partitions, the structure $\mathcal{T}_k^{\sqcup}(2)$ is especially relevant. For $k = 2$ it is isomorphic to the structure of levels of the Wagner hierarchy.



Picture 1: An initial segment of \mathcal{F}_2 .



Picture 2: An initial segment of \mathcal{F}_3 .



Picture 3: A fragment of $\mathcal{T}_3^{\omega}(2)$.

Theorem

1. *The quotient-posets of $(\mathcal{R}_k; \leq_{CA})$ and of $(\mathcal{R}_k; \leq_{DA})$ are isomorphic to the quotient-poset of $\mathcal{T}_k^{\sqcup}(2)$.*
2. *The relations \leq_{CA}, \leq_{DA} coincide on \mathcal{R}_k , the same holds for the relations \leq_{CS}, \leq_{DS} .*
3. *The relations $L(\mathcal{A}, c) \leq_{CA} L(\mathcal{A}, c)$ and $L(\mathcal{A}, c) \leq_{DA} L(\mathcal{A}, c)$ are decidable.*





- 1) Extending and modifying some operations of W. Wadge and A. Andretta on subsets of the Cantor space, we embed $\mathcal{T}_k^{\sqcup}(2)$ into $(\mathcal{R}_k; \leq_{CA})$ and $(\mathcal{R}_k; \leq_{DA})$ (an embedding is induced by $F \mapsto r(F)$).
- 2) We extend the author FH of sets [S98] to the FH of k -partitions over $(\Sigma_1^0 \cap \mathcal{R}, \Sigma_2^0 \cap \mathcal{R})$ in such a way that $r(F)$ is CA -complete in $\Sigma(F)$ and DA -complete in $\Sigma\mathcal{R}(F)$.
- 3) Relate to any Muller k -acceptor $\mathcal{A} = (\mathcal{A}, c)$ the structure $(C_{\mathcal{A}}; \leq_0, \leq_1, c)$ where $C_{\mathcal{A}}$ is the set of cycles of \mathcal{A} , $D \leq_0 E$ iff some state in D is reachable in the graph of the automaton \mathcal{A} from some state in E , and $D \leq_1 E$ iff $D \subseteq E$.






- 4) The structure $(C_{\mathcal{A}}; \leq_0, \leq_1, c)$ may be identified with some $P_{\mathcal{A}} \in \mathcal{P}_k(2)$.
- 5) Using the known facts [S98] that $(\Sigma_1^0 \cap \mathcal{R}, \Sigma_2^0 \cap \mathcal{R})$ have the reduction property conclude that $\Sigma\mathcal{R}(P_{\mathcal{A}}) = \Sigma\mathcal{R}(F_{\mathcal{A}})$ where $F_{\mathcal{A}} \in \mathcal{T}_k^{\sqcup}(2)$ is the natural unfolding of $P_{\mathcal{A}}$.
- 6) Check that $L\mathcal{A}$ is CA -complete in $\Sigma(F_{\mathcal{A}})$ and DA -complete in $\Sigma\mathcal{R}(F_{\mathcal{A}})$ and conclude that $L\mathcal{A} \equiv_{DA} r\mathcal{A}$.

So far, our results for ω -regular partitions generalized the corresponding results for ω -regular languages. Now we present a result that has completely different formulations for ω -regular languages and for ω -regular k -partitions for $k > 2$. Recall that *first-order theory* $FO(A)$ of a structure A of signature σ is the set of first-order sentences of signature σ which are true in A .

Theorem

For any $k \geq 3$, $FO(\mathcal{R}_k; \leq_{CA})$ is undecidable and, moreover, it is computably isomorphic to the first-order arithmetic $FO(\omega; +, \cdot)$.

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