

# EFFECTIVE FRACTAL DIMENSIONS

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## Lectures

1. Classical and Constructive Fractal Dimensions
2. Dimensions of Points in Euclidean Space
3. Finite-State Dimensions and Polynomial-Time Dimensions

# Topics for Today

## Classical fractal dimensions

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Effectivizing fractal dimensions

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Kolmogorov complexity characterizations of dimension



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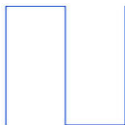
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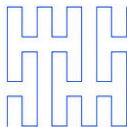
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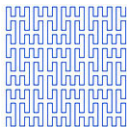
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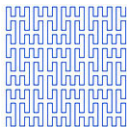
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Hausdorff, 1919: Rigorous formulation of dimension.



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- Open ball

$$B^\circ(x, r) = \{ y \in \mathcal{X} \mid \rho(x, y) < r \}.$$



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$$P^s(X) = \inf \left\{ \sum_{i=0}^{\infty} P_0^s(X_i) \mid X \subseteq \bigcup_{i=0}^{\infty} X_i \right\} \\ \text{optimize over all countable partitions of } X$$

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$H^s(\cdot)$  and  $P^s(\cdot)$  are outer measures.

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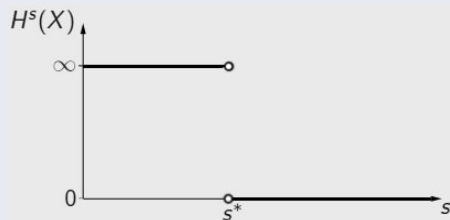
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Let  $\rho$  be a metric on  $\mathcal{X}$ , and let  $X \subseteq \mathcal{X}$ .

1. (Hausdorff 1919) The Hausdorff dimension of  $X$  with respect to  $\rho$  is

$$\dim^{(\rho)}(X) = \inf \{s \mid H^s(X) = 0\}.$$



$H^{s^*}(X)$  can be anything.

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2. (Tricot 1982, Sullivan 1984) The packing dimension of  $X$  with respect to  $\rho$  is

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# Fractal Dimension in Sequence Spaces

Let  $\Sigma$  be an alphabet with  $2 \leq |\Sigma| < \infty$ .

A (Borel) probability measure on  $\Sigma^\infty$  is a function  $\nu : \Sigma^* \rightarrow [0, 1]$  satisfying

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We restrict attention to probability measure that are strongly positive, meaning that there exists  $\delta > 0$  such that, for all  $w \in \Sigma^*$  and  $a \in \Sigma$ ,  $\nu(wa) \geq \delta \nu(w)$ .

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The metric induced by a strongly positive probability measure  $\nu$  on  $\Sigma^\infty$  is the function

$$\rho_\nu : \Sigma^\infty \times \Sigma^\infty \rightarrow [0, 1]$$

$$\rho_\nu(S, T) = \sup \{ \nu(w) \mid w \sqsubseteq S \text{ and } w \sqsubseteq T \}.$$

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$\dim^\nu(X)$  is also called the **Billingsley dimension** of  $X$  (Billingsley 1960).

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When  $\nu = \mu$ , we omit it from the terminology:

- The Hausdorff dimension of  $X$  is  $\dim_{\text{H}}(X) = \dim^\mu(X)$ .
- The packing dimension of  $X$  is  $\dim_{\text{P}}(X) = \text{Dim}^\mu(X)$ .

# Gale Characterizations

In a few minutes, we will define martingales, gales, and conditions for their success.

For the moment, martingales are strategies for betting on the successive symbols in a sequence  $S \in \Sigma^\infty$ , and one of these strategies succeeds on  $S$  if it makes an infinite amount of money betting on  $S$ .

Gales are generalized martingales that are no more powerful, but exhibit the martingales' success rates in a convenient form.

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Athreya, Hitchcock, Lutz, and Mayordomo 2007: Same for packing dimension.

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Schnorr, 1970's: Effective martingale success and success rates characterize certain types of randomness.

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Ryabko and Staiger, 1990's: Effective martingale success rates are related to Hausdorff dimension and Kolmogorov complexity.

Lutz, 2000: Martingale success rates characterize Hausdorff dimension. So used effective martingale success rates to define effective versions of Hausdorff dimension.

Athreya, Hitchcock, Lutz, and Mayordomo 2007:  
Same for packing dimension.

Lutz and Mayordomo 2008:  
Same for Billingsley dimensions.

## Definition

Let  $\nu$  be a probability measure on  $\Sigma^\infty$ , and let  $s \geq 0$ .

1. A  $\nu$ - $s$ -gale is a function  $d : \Sigma^* \rightarrow [0, \infty)$  that satisfies

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Observation (Lutz 2000)

$d$  is a  $\nu$ -s-gale  $\iff d'(w) = \nu(w)^{s-1}d(w)$  is a  $\nu$ -martingale.  
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# Gale Characterizations

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3. The success set of  $d$  is  $S^\infty[d] = \{S \mid d \text{ succeeds on } S\}$ .
4. The strong success set of  $d$  is  $S_{\text{str}}^\infty[d] = \{S \mid d \text{ succeeds strongly on } S\}$ .

## Theorem (Lutz & Mayordomo, 2008)

Let  $\nu$  be a strongly positive probability measure on  $\Sigma^\infty$ , and let  $X \subseteq \Sigma^\infty$ .

1. The Billingsley  $\nu$ -dimension of  $X$  is

$$\dim^\nu(X) = \inf \left\{ s \mid \begin{array}{l} \text{there is a } \nu\text{-s-gale } d \\ \text{such that } X \subseteq S^\infty[d] \end{array} \right\}.$$

2. The strong Billingsley  $\nu$ -dimension of  $X$  is

$$\text{Dim}^\nu(X) = \inf \left\{ s \mid \begin{array}{l} \text{there is a } \nu\text{-s-gale } d \\ \text{such that } X \subseteq S_{\text{str}}^\infty[d] \end{array} \right\}.$$

Recall

3. The Hausdorff dimension of  $X$  is

$$\dim_H(X) = \dim^\mu(X).$$

4. The packing dimension of  $X$  is

$$\dim_P(X) = \text{Dim}^\mu(X).$$

# Effective Fractal Dimensions

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Let  $\Delta$  be a resource bound, such as computable, constructive, poly-time, or finite-state.

We obtain  $\Delta$ -effective dimensions by requiring the gales in the gale characterizations to be  $\Delta$ -computable.

# Constructive Dimensions

A very important case is  $\Delta = \text{constructive}$ .

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## Definition

A  $\nu$ -s-gale  $d$  is constructive if it is lower semi-computable, i.e., if there is an exactly computable function  $\hat{d} : \Sigma^* \times \mathbb{N} \rightarrow \mathbb{Q}$  with the following two properties.

- For all  $w \in \Sigma^*$  and  $t \in \mathbb{N}$ ,  $\hat{d}(w, t) \leq \hat{d}(w, t + 1) < d(w)$ .
- For all  $w \in \Sigma^*$ ,  $\lim_{t \rightarrow \infty} \hat{d}(w, t) = d(w)$ .

# Constructive Dimensions

Definition (Lutz and Mayordomo 2008, aided by a result of Fenner 2002 )

Let  $\nu$  be a strongly positive probability measure on  $\Sigma^\infty$ , and let  $X \subseteq \Sigma^\infty$ .

1. The constructive  $\nu$ -dimension of  $X$  is

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2. The constructive strong  $\nu$ -dimension of  $X$  is

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We write

$$\text{cdim}(X) = \text{cdim}^\mu(X)$$

and

$$\text{cDim}(X) = \text{cDim}^\mu(X).$$

# Constructive Dimensions

A correspondence principle for an effective dimension is a theorem stating that, on sufficiently simple sets, the effective dimension coincides with its classical counterpart. (Terminology stolen from N. Bohr by Lutz.)

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Correspondence Principle for Constructive Dimension

Theorem ( Hitchcock 2002 )

*If  $X \subseteq \Sigma^\infty$  is any union (not necessarily effective) of computably closed (i.e.,  $\Pi_1^0$ ) sets then  $\text{cdim}(X) = \text{dim}_H(X)$ .*



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Correspondence Principle for Constructive Strong Dimension  
is false! (Conidis 2009)

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1. The  $\nu$ -dimension of  $S$  is  $\dim^\nu(S) = \text{cdim}^\nu(\{S\})$ .
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# Individual Sequences

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## Absolute Stability of Constructive Dimensions

Theorem ( [Lutz and Mayordomo 2008](#), extending [Lutz 2000](#) )

If  $\nu$  is a strongly positive, computable probability measure on  $\Sigma^\infty$ ,  
the, for all  $X \subseteq \Sigma^\infty$ ,

$$\text{cdim}^\nu(X) = \sup_{S \in X} \dim^\nu(S)$$

and

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(Contrast with countable stability of classical dimensions.)

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(Contrast with countable stability of classical dimensions.)

$\therefore$  Constructive dimensions are investigated in terms of individual sequences.

# Individual Sequences

In general,

$$0 \leq \dim_H(X) \leq \dim_P(X)$$
$$\quad \quad \quad \wedge \quad \quad \quad \wedge$$
$$\text{cdim}(X) \leq \text{cDim}(X) \leq 1.$$

## Definition (Martin-Löf '66, Schnorr '70)

A sequence  $R \in \mathcal{C}$  is random if no constructive martingale succeeds on  $R$ .

# Individual Sequences

If  $R$  is random (with respect to the uniform probability measure on  $\mathcal{C}$ ), then

$$\dim(R) = \text{Dim}(R) = 1.$$



# Individual Sequences

What if  $R$  is random with respect to some other probability measure on  $\mathcal{C}$ ?

# Individual Sequences

Fix  $\delta > 0$  and a bias sequence  $\vec{\beta} = (\beta_0, \beta_1, \beta_2, \dots)$  with each  $\beta_i \in [\delta, 1 - \delta]$ .

## Definition

$$\mathcal{H}(\beta) = \beta \log \frac{1}{\beta} + (1 - \beta) \log \frac{1}{1 - \beta} = \text{Shannon entropy.}$$

$$H_n(\vec{\beta}) = \frac{1}{n} \sum_{i=0}^{n-1} \mathcal{H}(\beta_i)$$

$$H^-(\vec{\beta}) = \liminf_{n \rightarrow \infty} H_n(\vec{\beta}) \quad \text{lower average entropy}$$

$$H^+(\vec{\beta}) = \limsup_{n \rightarrow \infty} H_n(\vec{\beta}) \quad \text{upper average entropy}$$

## Theorem (Athreya, Hitchcock, Lutz, & Mayordomo '07)

Let  $0 < \delta \leq \frac{1}{2}$ , and let  $\vec{\beta} = (\beta_0, \beta_1, \dots)$  be a computable bias sequence with each  $\beta_i \in [\delta, \frac{1}{2}]$ . For every  $\vec{\beta}$ -random sequence  $R$  we have

$$\dim(R) = H^-(\vec{\beta}), \quad \text{Dim}(R) = H^+(\vec{\beta}).$$

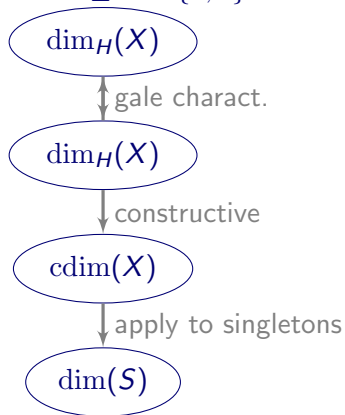
# Dimensions of finite strings

Our main task: Extend Hausdorff dimension to define  $\dim(x)$  for each  $x \in 0,1^*$ .

# Dimensions of finite strings

Our strategy:

For  $X \subseteq \mathcal{C} = \{0, 1\}^\infty$

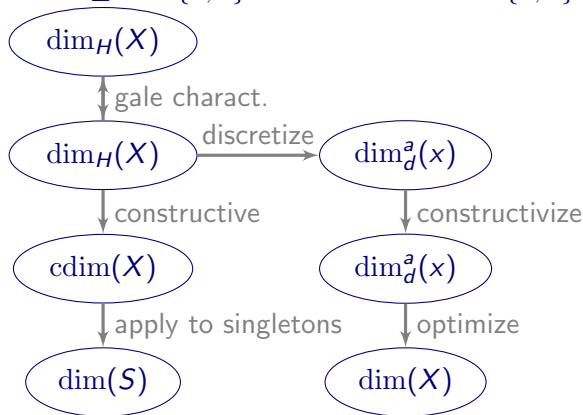


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# Dimensions of finite strings

Notation:  $\mathcal{T} = \underbrace{\{0,1\}^*}_{\text{finite strings}} \cup \underbrace{\{0,1\}^*\square}_{\text{finite strings with a blank symbol}}$

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Notation:  $\mathcal{T} = \underbrace{\{0,1\}^*}_{\text{prefixes thereof}} \cup \underbrace{\{0,1\}^*\square}_{\text{terminated binary strings}}$

## Definition

An s-termgale is a function  $d : \mathcal{T} \rightarrow [0, \infty)$  satisfying

$$d(\lambda) \leq 1$$

and

$$d(w) \geq 2^{-s} [d(w0) + d(w1) + d(w\Box)]$$

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Bets on the successive bits **and termination** of a finite string.

# Dimensions of finite strings

## Example

Define  $d : \mathcal{T} \rightarrow [0, \infty)$  by

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then  $d(w\Box) \gg d(\lambda)$ , even though  $d$  loses  $\frac{3}{4}$  of its money when the  $\Box$  appears.



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Trivial observation: If

$$2^{-s|x|}d(x) = 2^{-s'|x|}d'(x)$$



for all  $x \in \mathcal{T}$ , then  $d$  is an  $s$ -termgale  $\Leftrightarrow d'$  is an  $s'$ -termgale.

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$d$  is completely determined by any one of its elements.

# Dimensions of finite strings

## Definition

Let  $d$  be a termgale,  $a \in \mathbb{Z}^+$ , and  $w \in \{0, 1\}^*$ . The dimension of  $w$  relative to  $d$  at significance level  $a$  is

$$\dim_d^a(w) = \inf \left\{ s \mid d^{(s)}(w \square) > a \right\}.$$

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We have now discretized Hausdorff dimension. Constructivizing is easy:

## Definition

A termgale  $d$  is constructive if  $d^{(0)}$  is lower semicomputable.

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Now optimize:

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A constructive termgale  $\tilde{d}$  is optimal if for every constructive termgale  $d$  there is a constant  $\alpha > 0$  such that, for all  $s \in [0, \infty)$  and  $w \in \{0, 1\}^*$ ,

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## Theorem (Lutz 2003)

If  $\tilde{d}$  is an optimal constructive termgale, then, for every constructive termgale  $d$  and every  $a \in \mathbb{Z}^+$ , there is a constant  $\gamma \in [0, \infty)$  such that, for all  $w \in \{0, 1\}^*$ ,

$$\dim_{\tilde{d}}^a(w) \leq \dim_d(w) + \frac{\gamma}{1 + |w|}.$$

# Dimensions of finite strings

## Corollary

If  $d_1, d_2$  are optimal constructive termgales and  $a_1, a_2 \in \mathbb{Z}^+$ , then there is a constant  $\alpha \in [0, \infty)$  such that, for all  $w \in \{0, 1\}^*$ ,

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$\therefore$  It makes very little difference which optimal constructive termgale or which significance level we use.

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*Proof uses Levin's  $\mathbf{m}$ .*

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The dimension of a string  $w \in \{0, 1\}^*$  is

$$\dim(w) = \dim_{\mathbf{d}_{\square}}^1(w).$$

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*There is a constant  $c \in \mathbb{N}$  such that, for all  $x \in 0,1^*$ ,*

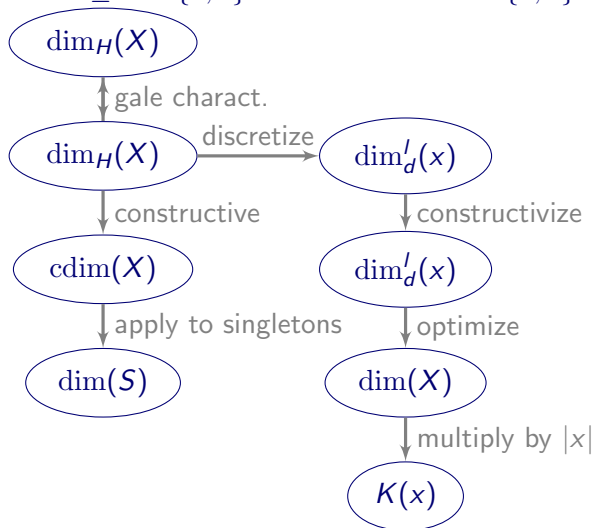
$$|K(x) - |x|\dim(x)| \leq c.$$

# Dimension and Kolmogorov Complexity

Our strategy:

For  $X \subseteq \mathcal{C} = \{0, 1\}^\infty$

For  $x \in \{0, 1\}^*$



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∴ Up to constant additive terms,

$$K(x) = \log \frac{1}{\mathbf{m}(x)} = |x| \dim(x).$$

Three fundamentally different approaches to information, when constructivized and optimized (after discretizing  $\dim_H$ ) lead to the same fundamental quantity,  $K(x)$ .



# Kolmogorov Complexity Characterizations

Our dimension characterization of  $K(\alpha)$  is one of several ways to prove the following.

Theorem (Lutz and Mayordomo 2008)

If  $\nu$  is a strongly positive, computable probability measure on  $\Sigma^\infty$ , then, for all  $S \in \Sigma^\infty$ ,

$$\dim^\nu(S) = \liminf_{m \rightarrow \infty} \frac{K(S[0..m-1])}{\mathcal{I}_\nu(S[0..m-1])}$$

and

$$\text{Dim}^\nu(S) = \limsup_{m \rightarrow \infty} \frac{K(S[0..m-1])}{\mathcal{I}_\nu(S[0..m-1])},$$

where

$K(w)$  = Kolmogorov complexity of  $w$  (any flavor)

and

$\mathcal{I}(w) = \text{Shannon } \nu\text{-self-information of } w = \log \frac{1}{\nu(w)}.$

Thank you!